1. Intro. This is a quick-and-dirty way to go from a slightly symbolic description of proposed mutual-exclusion algorithms to a corresponding set of clauses, so that I can use the clauses for bounded model checking.

In other words, I want to see whether the given concurrent algorithms can violate the mutex property by permitting simultaneous execution of two critical sections, or whether they can lead to livelock or starvation, in a given number of steps. To test this, I’ll see if certain extensions of the clauses are satisfiable.

First I have to describe the input language. Each step/state of an algorithm is given a name, which begins with an uppercase letter and has at most four characters. Every shared variable is also given a number, which begins with a lowercase letter and has at most two characters.

Only four elementary kinds of primitive operations are permitted at each step:

1) Compute non-critically, then optionally go to step $l$. (Here $l$ is a step name.)
2) Compute critically, then go to step $l$. (Likewise.)
3) Set $V \leftarrow v$, then goto $l$. (Here $V$ is a shared variable and $v$ is a constant.)
4) If $V = v$, goto $l$, else goto $l'$. (Likewise.)

These steps specify state transitions in an fairly obvious way; precise semantics will be explained later.

Here’s a simple example of possible input:

```
~ separate locks
A0 maybe goto A1
A1 a=1 goto A2
A2 if b=1 goto A2 else A3
A3 critical goto A4
A4 a=0 goto A0
B0 maybe goto B1
B1 b=1 goto B2
B2 if a=1 goto B2 else B3
B3 critical goto B4
B4 b=0 goto B0
```

The first line, which begins with `~", is simply a comment that will be passed to the output file. It is followed by steps of types 1, 3, 4, 2, 3, 1, 3, 4, 2, 3, respectively. The shared variables are $a$ and $b$. The concurrent occurrence of critical states should never occur.

(I do not claim that these programs solve the mutex problem; they simply provide an example.)

At present I assume that all step names begin with either A or B, and that all shared variables are Boolean. But those restrictions might well be lifted later, after I get some experience with this simpler scheme.
2. Here then is the basic outline of this program.

```c
#define maxsteps 100 /* at most this many steps */
#define bufsize 1024 /* must exceed the length of the longest input line */
#include <stdio.h>
#include <stdlib.h>
#include <string.h>

/* Type definitions */

typedef struct state {
    char name[4];    /* the name of this step and its successors */
    char var[2];     /* the shared variable */
    char val;        /* its value */
    char crit;       /* is this a critical step? */
} step;
```

This code is used in section 2.

3. (Process the command line 3) \(\equiv\)

```c
if (argc \neq 2 \lor sscanf(argv[1], "%d", &r) \neq 1) {
    fprintf(stderr, "Usage: sat-mutex <foo.dat> argv[0]);
    exit(-1);
}
if (r \leq 0) {
    fprintf(stderr, "Parameter r must be positive!\n"):
    exit(-2);
}
printf("sat-mutex %d\n", r);
```

This code is used in section 2.

4. Every non-comment line of input is recorded in an abbreviated form.

   (Type definitions 4) \(\equiv\)

```c
typedef struct state_struct {
    char name[4], lab[4], elab[4];    /* the name of this step and its successors */
    char var[2];    /* the shared variable */
    char val;       /* its value */
    char crit;      /* is this a critical step? */
} step;
```

This code is used in section 2.
5. I don’t attempt to provide much syntactic sugar for the user (since I expect to be the only user). If I need something fancier, I’ll probably write a preprocessor to convert fancy output into the primitive form that is understood by this program.

(Parse the input into the state table 5) \[\]

\[
\text{for } (m = n = ma = mb = 0; ; ) \{
\text{if } (!\text{fgets(buf, bufsize, stdin)}) \text{ break;)
\text{if } (buf[0] \equiv \text{""}) \text{ printf("} %s \text{"}, buf);
\text{else } \{
\text{char } \ast\text{curp} = \text{buf};
\text{if } (m \geq \text{maxsteps}) \{
\text{fprintf(stderr, "Recompile me, I only have room for \#d steps!\n", maxsteps);}
\text{exit}(-666);
\}
\langle \text{Scan the name field 6};
\text{if } (\text{strcmp(curp, } \text{"maybe", } 6) \equiv 0) \{ \text{Scan a maybe step 7}
\text{else if } (\text{strcmp(curp, } \text{"critical", } 9) \equiv 0) \{ \text{Scan a critical step 8}
\text{else if } (\text{strcmp(curp, } \text{"if", } 3) \equiv 0) \{ \text{Scan an if step 10}
\text{else } (\text{Scan an assignment step 14});
\text{m}++;\}
\}\langle \text{Check for missing steps 15};
\text{if } (\text{state[astep[0]].crit + state[bstep[0]].crit} > 1) \{
\text{fprintf(stderr, "Both processes are initially in critical sections!\n");}
\text{exit}(-555);
\}
\text{fprintf(stderr, "\(\#d\%d\_steps\_with\_%d\_shared\_variables\_successfully\_input\)\n", ma, mb, n);}
\text{This code is used in section 2.}
\]

6. \#define \textbf{abrt}(m, t)
\{ \text{fprintf(stderr, "Oops, \#s\n", m, buf); exit(t); } \}

\langle \text{Scan the name field 6}; \equiv
\text{for } (j = 0; \ast\text{curp} \&\& \ast\text{curp} \neq \text{"", } \text{\&}\text{\&} \ast\text{curp} \neq \text{"n"); } j++, \text{curp}++)
\text{if } (j < 4) \text{state[\text{m}.name[j]} = \ast\text{curp};
\text{if } (j > 4) \text{abrt("the name is too long", } -10);\text{if } (\text{state[\text{m}.name[0]} < \text{A} \vee \text{state[\text{m}.name[0]} > \text{B})
\text{abrt("the step name must begin with A or B", } -11);\text{for } (j = 0; j < m; j++)
\text{if } (\text{strcmp(state[j].name, state[\text{m}.name, 4]} \equiv 0) \text{abrt("that name has already been used", } -12);\text{if } (\text{state[\text{m}.name[0]} \equiv \text{A}) \text{ asstep[ma++] } = m;
\text{else } \text{bstep[mb++] } = m;
\text{if } (\ast\text{curp} \neq \text{"n"); abrt("step is incomplete", } -13);\text{This code is used in section 5.}
\]

7. \langle \text{Scan a maybe step 7}; \equiv
\{ \text{curp } += 5;
\langle \text{Scan the lab field 9};
\text{if } (\ast\text{curp} \neq \text{"n"); abrt("maybe step ends badly", } -14);\text{This code is used in section 5.}
\]
8. \(\text{(Scan a critical step 8)}\) \(\equiv\)
\[
\begin{align*}
&\text{curp} += 8; \\
&\text{state}[m].\text{crit} = 1; \\
&(\text{Scan the lab field 9}); \\
&\text{if } (*\text{curp} \neq ^'\text{\n}') \text{ \text{\text{abrt}}("critical step ends badly", -15);} \\
&\}
\]
This code is used in section 5.

9. \(\text{(Scan the lab field 9)}\) \(\equiv\)
\[
\begin{align*}
&\text{if } (\text{strncpy(curp, "\_goto\", 6) \neq 0}) \text{ \text{\text{abrt}}("missing goto", -16);} \\
&\text{curp} += 6; \\
&\text{for } (j = 0; *\text{curp} \land *\text{curp} \neq ^'_{j} \land *\text{curp} \neq ^'\text{\n}'; j++, \text{curp}++) \\
&\quad \text{if } (j < 4) \text{ \text{\text{state}}[m].\text{lab}[j] = *\text{curp};} \\
&\quad \text{if } (j > 4) \text{ \text{\text{abrt}}("the label is too long", -17);} \\
&\}
\]
This code is used in sections 7, 8, 10, and 14.

10. \(\text{(Scan an if step 10)}\) \(\equiv\)
\[
\begin{align*}
&\text{curp} += 3; \\
&(\text{Scan the var field 11}); \\
&\text{if } (*\text{curp} ++ \neq ^'=' ) \text{ \text{\text{abrt}}("missing '=' in an if step", -18);} \\
&(\text{Scan the val field 12}); \\
&(\text{Scan the lab field 9}); \\
&(\text{Scan the elab field 13}); \\
&\text{if } (*\text{curp} \neq ^'\text{\n}') \text{ \text{\text{abrt}}("that if step ends badly", -19);} \\
&\}
\]
This code is used in section 5.

11. \(\text{(Scan the var field 11)}\) \(\equiv\)
\[
\begin{align*}
&\text{for } (j = 0; *\text{curp} \land *\text{curp} \neq ^'_{j} \land *\text{curp} \neq ^'\text{\n}'; j++, \text{curp}++) \\
&\quad \text{if } (j < 2) \text{ \text{\text{state}}[m].\text{var}[j] = \text{\text{vars}}[n][j] = *\text{curp};} \\
&\quad \text{if } (j > 2) \text{ \text{\text{abrt}}("a variable name is too long", -20);} \\
&\quad \text{if } (\text{\text{state}}[m].\text{var}[0] < ^'a' \lor \text{\text{state}}[m].\text{var}[0] > ^'z') \\
&\quad \quad \text{abrt("a variable name must begin with a lowercase letter", -21);} \\
&\quad \text{for } (j = 0; j < n; j++) \\
&\quad \quad \text{if } (\text{\text{strncpy(\text{vars}[j]}, \text{\text{state}}[m].\text{var}, 2) \equiv 0}) \text{ \text{\text{break}};} \\
&\quad \text{if } (j \equiv n) \text{ \text{\text{\text{n}}++};} \\
&\quad \text{else } \text{\text{vars}}[n][1] = 0; \\
&\}
\]
This code is used in sections 10 and 14.

12. \(\text{(Scan the val field 12)}\) \(\equiv\)
\[
\begin{align*}
&\text{if } (*\text{curp} < ^'0' \lor *\text{curp} > ^'1') \text{ \text{\text{abrt}}("the value must be 0 or 1", -22);} \\
&\text{\text{state}}[m].\text{val} = *\text{curp} ++ - ^'0'; \\
&\}
\]
This code is used in sections 10 and 14.
13. (Scan the \texttt{elab} field 13) \equiv
   \begin{align*}
   &\text{if} (\text{strncmp}(\text{curp}, \_\text{else}_n, 6) \neq 0) \text{ abrt}(\text{"missing}_n\_\text{else"}, -23); \\
   &\text{curp} += 6; \\
   &\text{for } (j = 0; *\text{curp} \land *\text{curp} \neq \text{\'\n\'}; j++, \text{curp}++) \\
   &\text{if } (j < 4) \text{ state}[m].\text{elab}[j] = *\text{curp}; \\
   &\text{if } (j > 4) \text{ abrt}(\text{"the}_n\_\text{else}_n\_\text{label}_n\_is\_too\_long"}, -24); \\
   \end{align*}
This code is used in section 10.

14. (Scan an assignment step 14) \equiv
   \begin{align*}
   &\text{(Scan the } \texttt{var} \text{ field 11);} \\
   &\text{if } (*\text{curp}++ \neq \text{\'=\'}) \text{ abrt}(\text{"missing}_n\_\text{else}_n\_\text{in}_n\_\text{an}_n\_\text{assignment}_n\_\text{step"}, -25); \\
   &\text{(Scan the } \texttt{val} \text{ field 12);} \\
   &\text{(Scan the } \texttt{lab} \text{ field 9);} \\
   &\text{if } (*\text{curp} \neq \text{\'\n\'}) \text{ abrt}(\text{"assignment}_n\_\text{step}_n\_\text{ends}_n\_badly"}, -26); \\
   \end{align*}
This code is used in section 5.

15. (Check for missing steps 15) \equiv
   \begin{align*}
   &\text{if } (ma \equiv 0) \{ \\
   &\quad \text{fprintf}(\text{stderr}, \text{"There}_n\_\text{are}_n\_no\_steps}_n\_\text{for}_n\_\text{process}_n\_A\_!\_\text{n"}); \\
   &\quad \text{exit}(\neg 99); \\
   &\} \\
   &\text{if } (mb \equiv 0) \{ \\
   &\quad \text{fprintf}(\text{stderr}, \text{"There}_n\_\text{are}_n\_no\_steps}_n\_\text{for}_n\_\text{process}_n\_B\_!\_\text{n"}); \\
   &\quad \text{exit}(\neg 98); \\
   &\} \\
   &\text{for } (k = t = 0; k < m; k++) \{ \\
   &\quad \text{if } (\text{state}[k].\text{lab}[0]) \{ \\
   &\quad\quad \text{for } (j = 0; j < m; j++) \\
   &\quad\quad\quad \text{if } (\text{strncmp}(\text{state}[j].\text{name}, \text{state}[k].\text{lab}, 4) \equiv 0) \text{ break;} \\
   &\quad\quad\quad \text{if } (j \equiv m) \{ \\
   &\quad\quad\quad\quad \text{fprintf}(\text{stderr}, \text{"Missing}_n\_\text{step}_n\_X\_!\_\text{n"}, \text{state}[k].\text{lab}); \\
   &\quad\quad\quad\quad t++; \\
   &\quad\quad\} \\
   &\quad\} \\
   &\quad \text{if } (\text{state}[k].\text{elab}[0]) \{ \\
   &\quad\quad \text{for } (j = 0; j < m; j++) \\
   &\quad\quad\quad \text{if } (\text{strncmp}(\text{state}[j].\text{name}, \text{state}[k].\text{elab}, 4) \equiv 0) \text{ break;} \\
   &\quad\quad\quad \text{if } (j \equiv m) \{ \\
   &\quad\quad\quad\quad \text{fprintf}(\text{stderr}, \text{"Missing}_n\_\text{step}_n\_X\_!\_\text{n"}, \text{state}[k].\text{elab}); \\
   &\quad\quad\quad\quad t++; \\
   &\quad\quad\} \\
   &\quad\} \\
   &\quad \text{if } (t) \text{ exit}(-30); \\
   \end{align*}
This code is used in section 5.
16. The generated clauses involve variables like ‘2A1’, meaning that process A is in state A1 at time 2; also variables like ‘3b’, meaning that shared variable b is 1 (true) at time 3; also variables like ‘1@’, meaning that process A took a turn at time 1. (The negations of these variables, namely ~2A1, ~3b, ~1@, mean respectively that A is not in state A1 at time 2, b is 0 (false) at time 3, and process B took a turn at time 1.) At time 0, all shared variables are 0 and each process is in its first-mentioned state.

(Generate the initial clauses 16) ≡ 
{ 
  for (j = 0; j < n; j++) printf("000%.2s\n", vars[j]);
  printf("000%.4s\n", state[astep[0]].name);
  for (j = 1; j < ma; j++) printf("000%.4s\n", state[astep[j]].name);
  printf("000%.4s\n", state[bstep[0]].name);
  for (j = 1; j < mb; j++) printf("000%.4s\n", state[bstep[j]].name);
}
This code is used in section 2.

17. Speaking of turns reminds me that I promised to define precise semantics.

At each time t one of the processes, chosen nondeterministically, is granted permission to take a turn, which means intuitively that it performs the step corresponding to its current state. We say that the selected process is “bumped.”

Every process is in a unique state at time t. The state of a process remains the same at time t + 1 if it’s not bumped. But if it’s bumped, the next state is (1) either the same or lab, nondeterministically, after a maybe step; (2) lab after a critical step or an assignment step; (2) either lab or elab after an if step, depending on whether or not the shared variable has the specified value.

The value of a shared variable at time t + 1 is the same as the value that it had at time t, unless the bumped process assigned another value to it. In particular, if two processes are trying to change the same shared variable, the bumped process changes it first.

When the bumped process executes an if statement at the same time as another process is trying to write the same variable, the other process does not influence the result of the if; the change it wants to make will have to wait. [This rule means that weaker algorithms can get by, but they need stronger (and presumably more expensive and/or slower) hardware support. I’m using this rule in all the early examples of mutex in TAOCP, because it is easier to explain; the harder rule can be considered later, after algorithms pass this simpler criterion.]

(Generate the transitions from time t to time t + 1 17) ≡ 
{ 
  (Generate clauses to forbid nonunique states for A at time t + 1 18);
  (Generate clauses to forbid nonunique states for B at time t + 1 19);
  (Generate the state transition clauses when A is bumped 20);
  (Generate the state transition clauses when B is bumped 22);
  (Generate the variable transition clauses 24);
}
This code is used in section 2.
18. I introduce auxiliary variables here, using Heule’s exclusion clauses, so that we don’t have quadratic blowup when the programs are large.

```c
#define printprevA()
    if (j) printf("%03d_%d", t + 1, i - 1);
    else printf("~%03d_%d", t + 1, state[astep[k - 1]].name);
(Generate clauses to forbid nonunique states for A at time t + 1) \equiv
k = ma;
if (k > 1) {
    i = j = 0;
    if (k == 2) printf("~%03d_%d", t + 1, state[astep[0]].name, t + 1, state[astep[1]].name);
while (k > 4) {  
    printprevA();
    printf("~%03d_%d", t + 1, state[astep[k - 2]].name);
    printprevA();
    printf("~%03d_%d", t + 1, state[astep[k - 3]].name);
    printprevA();
    printf("~%03d_%d", t + 1, i);
    printprevA();
    printf("~%03d_%d", t + 1, state[astep[k - 2]].name, t + 1, state[astep[k - 3]].name);
    printprevA();
    printf("~%03d_%d", t + 1, i);
    printprevA();
    printf("~%03d_%d", t + 1, state[astep[k - 2]].name, t + 1, state[astep[k - 3]].name);
    i++, j = 1, k = 2;
}
printprevA();
printf("~%03d_%d", t + 1, state[astep[k - 2]].name);
printprevA();
printf("~%03d_%d", t + 1, state[astep[k - 3]].name);
printf("~%03d_%d", t + 1, state[astep[k - 2]].name, t + 1, state[astep[k - 3]].name);
if (k > 3) {  
    printprevA();
    printf("~%03d_%d", t + 1, state[astep[k - 4]].name);
    printprevA();
    printf("~%03d_%d", t + 1, state[astep[k - 4]].name, t + 1, state[astep[k - 4]].name);
    printprevA();
    printf("~%03d_%d", t + 1, state[astep[k - 4]].name, t + 1, state[astep[k - 4]].name);
}
}
```

This code is used in section 17.
19. #define printprevB()
    if (j) printf("%03d_B%d", t + 1, i - 1);
    else printf("%03d%.4s", t + 1, state[bstep[k-1]].name);

(Generate clauses to forbid nonunique states for B at time t + 1)
\[ k = mb; \]

if (k > 1) {
    i = j = 0;
    if (k ≥ 2) printf("%03d%.4s", t + 1, state[bstep[0]].name, t + 1, state[bstep[1]].name);
    while (k > 4) {
        printprevB();
        printf("%03d%.4s", t + 1, state[bstep[k - 2]].name);
        printprevB();
        printf("%03d%.4s", t + 1, state[bstep[k - 3]].name);
        printprevB();
        printf("%03d_B%d", t + 1, i);
        printf("%03d%.4s", t + 1, state[bstep[k - 2]].name, t + 1, state[bstep[k - 3]].name);
        printf("%03d%.4s", t + 1, state[bstep[k - 2]].name, t + 1, state[bstep[k - 3]].name);
        i++, j = 1, k = 2;
    }
    printprevB();
    printf("%03d%.4s", t + 1, state[bstep[k - 2]].name);
    printprevB();
    printf("%03d%.4s", t + 1, state[bstep[k - 3]].name);
    printf("%03d%.4s", t + 1, state[bstep[k - 2]].name, t + 1, state[bstep[k - 3]].name);
    if (k > 3) {
        printprevB();
        printf("%03d%.4s", t + 1, state[bstep[k - 4]].name);
        printf("%03d%.4s", t + 1, state[bstep[k - 4]].name);
        printf("%03d%.4s", t + 1, state[bstep[k - 4]].name);
    }
}

This code is used in section 17.

20. #define tprime (t + 1)

(Generate the state transition clauses when A is bumped)
\[ k = mb; k++; \]

for (k = 0; k < mb; k++) {
    printf("%03d%.4s", t, t, state[astep[k]].name, tprime, state[astep[k]].name);
    if (state[astep[k]].var[0] == 0) {
        if (state[astep[k]].crit == 0) /* a maybe step */
            printf("%03d%.4s", t, t, state[astep[k]].name, tprime, state[astep[k]].lab);
        else printf("%03d%.4s", t, t, state[astep[k]].name, tprime, state[astep[k]].lab);
    }
    else if (state[astep[k]].lab[0] == 0) { /* an assignment step */
        printf("%03d%.4s", t, t, state[astep[k]].name, tprime, state[astep[k]].lab);
    }
    else (Generate clauses for when A is bumped in an if step 21);
\[\text{Generate clauses for when A is bumped in an if step 21}\] \[\underline{\text{21.}} \quad \langle \text{Generate clauses for when A is bumped in an if step 21} \rangle \equiv \]

\[
\begin{array}{c}
\text{printf ("%03d\%s\%s\%s\%s\%s\%s\%s\%s\\n", t, t, state[astep[k]].name);} \\
\text{printf ("%s%03d\%s\%s\%s\%s\%s\%s\%s\\n", state[astep[k]].val ? "-" : ",", t, state[astep[k]].var, tprime, state[astep[k]].lab);} \\
\text{printf ("%03d\%s\%s\%s\%s\%s\%s\%s\%s\\n", t, t, state[astep[k]].name);} \\
\text{printf ("%s%03d\%s\%s\%s\%s\%s\%s\%s\\n", state[astep[k]].val ? "-" : ",", t, state[astep[k]].var, tprime, state[astep[k]].lab);}
\end{array}
\]

This code is used in section 20.

\[\text{Generate the state transition clauses when B is bumped 22}\] \[\underline{\text{22.}} \quad \langle \text{Generate the state transition clauses when B is bumped} \rangle \equiv \]

\[
\begin{array}{c}
\text{for (k = 0; k < mb; k++)} \{ \\
\text{printf ("%03d\%s\%s\%s\%s\%s\%s\%s\%s\\n", t, t, state[bstep[k]].name, tprime, state[bstep[k]].name);} \\
\text{if (state[bstep[k]].var[0] \equiv 0)} \{ \\
\text{if (state[bstep[k]].crit \equiv 0)} \quad /* \text{a maybe step} */ \\
\text{printf ("%03d\%s\%s\%s\%s\%s\%s\%s\%s\\n", t, t, state[bstep[k]].name, tprime, state[bstep[k]].name, tprime, state[bstep[k]].lab);} \\
\text{else printf ("%03d\%s\%s\%s\%s\%s\%s\%s\%s\\n", t, t, state[bstep[k]].name, tprime, state[bstep[k]].lab);} \\
\text{/* a critical step */}
\} \text{else if (state[bstep[k]].lab[0] \equiv 0)} \quad /* \text{an assignment step} */ \\
\text{printf ("%03d\%s\%s\%s\%s\%s\%s\%s\%s\\n", t, t, state[bstep[k]].name, tprime, state[bstep[k]].lab);} \\
\} \text{else (Generate clauses for when B is bumped in an if step 23);} \\
\end{array}
\]

This code is used in section 17.

\[\text{Generate clauses for when B is bumped in an if step 23}\] \[\underline{\text{23.}} \quad \langle \text{Generate clauses for when B is bumped in an if step 23} \rangle \equiv \]

\[
\begin{array}{c}
\text{printf ("%03d\%s\%s\%s\%s\%s\%s\%s\%s\\n", t, t, state[bstep[k]].name);} \\
\text{printf ("%s%03d\%s\%s\%s\%s\%s\%s\%s\\n", state[bstep[k]].val ? "-" : ",", t, state[bstep[k]].var, tprime, state[bstep[k]].lab);} \\
\text{printf ("%03d\%s\%s\%s\%s\%s\%s\%s\%s\\n", t, t, state[bstep[k]].name);} \\
\text{printf ("%s%03d\%s\%s\%s\%s\%s\%s\%s\\n", state[bstep[k]].val ? "-" : ",", t, state[bstep[k]].var, tprime, state[bstep[k]].lab);}
\end{array}
\]

This code is used in section 22.
24. \( \text{(Generate the variable transition clauses 24)} \equiv \)
\[ \text{for } (k = 0; k < n; k++) \{ \quad \text{/* first consider all cases where the value changes */} \]
\[ \text{for } (j = 0; j < m; j++) \]
\[ \text{if } (\text{strcmp}(\text{state}[j].\text{var}, \text{vars}[k], 2) \equiv 0 \land \text{state}[j].\text{elab}[0] \equiv 0) \]
\[ \text{printf ("%s%03d%\%03d\%, 4s, %s%03d\%, 2s\n", state[j].name[0] \equiv 'A' ? "\%s": "", t, state[j].name, state[j].val \equiv 0 ? "-": "", tprime, state[j].var);} \]
\[ \text{/* now consider all cases where the value doesn't change */} \]
\[ \text{printf ("C0%03d%", t, t, vars[k]);} \quad \text{/* A bumped and val is 0 */} \]
\[ \text{for } (j = 0; j < m; j++) \]
\[ \text{if } (\text{strcmp}(\text{state}[j].\text{var}, \text{vars}[k], 2) \equiv 0 \land \text{state}[j].\text{elab}[0] \equiv 0 \land \text{state}[j].\text{name}[0] \equiv 'A') \]
\[ \text{printf ("%03d\%s\%03d\%, 4s", t, state[j].name);} \quad \text{/* not changed by A */} \]
\[ \text{printf ("-%03d\%, 2s\n", tprime, vars[k]);} \quad \text{/* it stays 0 */} \]
\[ \text{printf ("-%03d\%, 2s", t, t, vars[k]);} \quad \text{/* B bumped and val is 0 */} \]
\[ \text{for } (j = 0; j < m; j++) \]
\[ \text{if } (\text{strcmp}(\text{state}[j].\text{var}, \text{vars}[k], 2) \equiv 0 \land \text{state}[j].\text{elab}[0] \equiv 0 \land \text{state}[j].\text{name}[0] \equiv 'B') \]
\[ \text{printf ("%03d\%s\%03d\%, 4s", t, state[j].name);} \quad \text{/* not changed by B */} \]
\[ \text{printf ("-%03d\%, 2s\n", tprime, vars[k]);} \quad \text{/* it stays 0 */} \]
\[ \text{printf ("-%03d\%, 2s", t, t, vars[k]);} \quad \text{/* A bumped and val is 1 */} \]
\[ \text{for } (j = 0; j < m; j++) \]
\[ \text{if } (\text{strcmp}(\text{state}[j].\text{var}, \text{vars}[k], 2) \equiv 0 \land \text{state}[j].\text{elab}[0] \equiv 0 \land \text{state}[j].\text{name}[0] \equiv 'A') \]
\[ \text{printf ("%03d\%s\%03d\%, 4s", t, state[j].name);} \quad \text{/* not changed by A */} \]
\[ \text{printf ("-%03d\%, 2s\n", tprime, vars[k]);} \quad \text{/* it stays 1 */} \]
\[ \text{printf ("-%03d\%, 2s", t, t, vars[k]);} \quad \text{/* B bumped and val is 1 */} \]
\[ \text{for } (j = 0; j < m; j++) \]
\[ \text{if } (\text{strcmp}(\text{state}[j].\text{var}, \text{vars}[k], 2) \equiv 0 \land \text{state}[j].\text{elab}[0] \equiv 0 \land \text{state}[j].\text{name}[0] \equiv 'B') \]
\[ \text{printf ("%03d\%s\%03d\%, 4s", t, state[j].name);} \quad \text{/* not changed by B */} \]
\[ \text{printf ("-%03d\%, 2s\n", tprime, vars[k]);} \quad \text{/* it stays 1 */} \]
\[ \text{printf ("-%03d\%, 2s", t, t, vars[k]);} \quad \text{/* B bumped and val is 1 */} \]
\} \]

This code is used in section 17.

25. The different ways of going jointly critical are \( C0, C1, \text{etc.} \)
(Generate clauses to force concurrent critical sections at time \( r \) 25) \( \equiv \)
\[ \text{for } (i = j = 0; j < ma; j++) \]
\[ \text{if } (\text{state[astep[j]].crit}) \{ \]
\[ \text{for } (k = 0; k < mb; k++) \]
\[ \text{if } (\text{state[bstep[k]].crit}) \{ \]
\[ \text{printf ("C%d\%03d\%, 4s\n", i, r, state[astep[j]].name);} \]
\[ \text{printf ("C%d\%03d\%, 4s\n", i, r, state[bstep[k]].name);} \]
\[ \text{i++;} \]
\[ \} \]
\[ \text{for } (j = 0; j < i; j++) \text{printf("C%d", j);} \]
\[ \text{printf ("\n");} \]

This code is used in section 2.
26. Index.

\textit{abrt}: 6, 7, 8, 9, 10, 11, 12, 13, 14.
\textit{argc}: 2, 3.
\textit{argv}: 2, 3.
\textit{astep}: 2, 5, 6, 16, 18, 20, 21, 25.
\textit{bstep}: 2, 5, 6, 16, 19, 22, 23, 25.
\textit{buf}: 2, 5, 6.
\textit{bufsize}: 2, 5.
\textit{crit}: 4, 5, 8, 20, 22, 25.
\textit{carp}: 5, 6, 7, 8, 9, 10, 11, 12, 13, 14.
\textit{elab}: 4, 13, 15, 17, 20, 21, 22, 23, 24.
\textit{exit}: 3, 5, 6, 15.
\textit{fgets}: 5.
\textit{fprintf}: 3, 5, 6, 15.
\textit{i}: 2.
\textit{j}: 2.
\textit{k}: 2.
\textit{lab}: 4, 9, 15, 17, 20, 21, 22, 23.
\textit{m}: 2.
\textit{ma}: 2, 5, 6, 15, 16, 18, 20, 25.
\textit{main}: 2.
\textit{maxsteps}: 2, 5.
\textit{mb}: 2, 5, 6, 15, 16, 19, 22, 25.
\textit{n}: 2.
\textit{name}: 4, 6, 15, 16, 18, 19, 20, 21, 22, 23, 24, 25.
\textit{printf}: 3, 5, 16, 18, 19, 20, 21, 22, 23, 24, 25.
\textit{printfB}: 3, 18.
\textit{printprevB}: 19.
\textit{r}: 2.
\textit{sscanf}: 3.
\textit{state}: 2, 5, 6, 8, 9, 11, 12, 13, 15, 16, 18, 19, 20, 21, 22, 23, 24, 25.
\textit{state struct}: 4.
\textit{stderr}: 3, 5, 6, 15.
\textit{stdin}: 2, 5.
\textit{step}: 2, 4.
\textit{strncpy}: 5, 6, 9, 11, 13, 15, 24.
\textit{t}: 2.
\textit{tprime}: 20, 21, 22, 23, 24.
\textit{val}: 4, 12, 21, 23, 24.
\textit{var}: 4, 11, 20, 21, 22, 23, 24.
\textit{vars}: 2, 11, 16, 24.
(Check for missing steps 15) Used in section 5.
(Generate clauses for when A is bumped in an if step 21) Used in section 20.
(Generate clauses for when B is bumped in an if step 23) Used in section 22.
(Generate clauses to forbid nonunique states for A at time $t + 1$) Used in section 17.
(Generate clauses to forbid nonunique states for B at time $t + 1$) Used in section 17.
(Generate clauses to force concurrent critical sections at time $r$) Used in section 2.
(Generate the initial clauses 16) Used in section 2.
(Generate the state transition clauses when A is bumped 20) Used in section 17.
(Generate the state transition clauses when B is bumped 22) Used in section 17.
(Generate the transitions from time $t$ to time $t + 1$) Used in section 2.
(Generate the variable transition clauses 24) Used in section 17.
(Parse the input into the state table 5) Used in section 2.
(Process the command line 3) Used in section 2.
(Scan a critical step 8) Used in section 5.
(Scan a maybe step 7) Used in section 5.
(Scan an if step 10) Used in section 5.
(Scan an assignment step 14) Used in section 5.
(Scan the elab field 13) Used in section 10.
(Scan the lab field 9) Used in sections 7, 8, 10, and 14.
(Scan the name field 6) Used in section 5.
(Scan the val field 12) Used in sections 10 and 14.
(Scan the var field 11) Used in sections 10 and 14.
(Type definitions 4) Used in section 2.
SAT-MUTEX

Section | Page
---|---
Intro | 1
Index | 26

1 | 1
11 | 11